## CSE211: Compiler Design

 Oct. 6, 2020- Topic: Parsing Overview

- Questions:
- What is parsing?
- Have you used Regular Expressions before?
- How do you parse Regular Expressions? What about Context-free Grammars?


## Announcements:

- Moving Homework due dates back one week (more time to work on homework after module is finished)
- Notes will include a reference to EAC
- Link to reserve is up
- How to watch YuJa recordings


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- How do you parse Regular Expressions? What about Context-free Grammars?


## Parsing is the first step in a compiler

- How do we parse language?


## Parsing is the first step in a compiler

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The dog ran across the park

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## ARTICLE NOUN VERB PREPOSITION ARTICLE NOUN

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ARTICLE NOUN VERB PREPOSITION ARTICLE NOUN

A Simple Language

- ARTICLE
- NOUN
- VERB
- ADJECTIVE


## A Simple Language

- ARTICLE $=$ \{The, A, My, Your\}
- NOUN = \{Dog, Car, Computer\}
- VERB $=\{$ Ran, Crashed, Accelerated $\}$
- ADJECTIVE $=\{$ Purple, Spotted, Microsoft $\}$


## A Simple Language

- ARTICLE $=$ \{The, A, My, Your\}
- NOUN = \{Dog, Car, Computer\}
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## ARTICLE NOUN VERB

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## ARTICLE ADJECTIVE? NOUN VERB

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## ARTICLE ADJECTIVE? NOUN VERB

My Microsoft Computer Crashed

## A Simple Language

- ARTICLE $=$ \{The, A, My, Your\}
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## ARTICLE ADJECTIVE? NOUN VERB

The Purple Dog Crashed

## Goals in this module

- Understand the architecture of a modern parser (tokenizing and parsing)
- Understand the language of tokens (regular expressions) and parsers (context-free grammars)
- How to design CFG production rules so avoid ambiguity and set precedence and associativity.
- Learn how to parse using a classic parser generator (Lex and Yacc) for a simple programming language


## Goals in this module

- We will NOT discuss parsing algorithms for CFGs. It is a deep dark hole. If you are interested, you can do this for a paper assignment.
- This module should provide you with the background to implement simple compilers. It will make you very popular with future colleagues who are scared of compilers.
- These topics are typically covered in more depth in an undergrad course (e.g. formal properties of regular expressions, parsing algorithms).

High-level parser

## Parser

## High-level parser



A string

## High-level parser



## High-level parser

The input string is in the language $L$


High-level parser of compilation


## High-level parser

continue to the rest of compilation


What happens here?

## ARTICLE ADJECTIVE? NOUN VERB

## High-level parser

continue to the rest of compilation


What happens

## ARTICLE ADJECTIVE? NOUN VERB

 The Purple Dog Crashed
## Parser architecture

## Parser



## Parser architecture

## Parser



## Parser architecture

## Parser



## Scanner

- List of tokens:
- e.g. \{NOUN, ARTICLE, ADJECTIVE, VERB\}


## Scanner

My Microsoft Computer Crashed

## Scanner

My Microsoft Computer Crashed


Scanner

(ARTICLE, my) (ADJECTIVE, Microsoft) (NOUN, Computer) (VERB, Crashed)

## Scanner

## My Microsoft Computer Crashed



## Scanner

(ARTICLE, my) (ADJECTIVE, Microsoft) (NOUN, Computer) (VERB, Crashed)
Lexeme: (TOKEN, value)

## Scanner

- Lets write tokens for arithmetic expression:

$$
(5+4) * 3
$$

## Scanner

- Lets write tokens for arithmetic expression:
(LPAREN, '(') (NUMBER, 5)
(PLUS, +)
(NUMBER, 4)
(RPAREN, ' ')')
(TIMES, *)
(NUMBER, 3)

$$
(5+4) * 3
$$

## Scanner

- Lets write tokens for arithmetic expression:
(LPAREN, '(') (NUMBER, 5)
(PLUS, +)
(NUMBER, 4)
(RPAREN, ' ')')
(TIMES, *)
(NUMBER, 3)

```
(5+4)*3
(LPAREN, '(')
(NUMBER,5)
(OP, +)
(NUMBER, 4)
(RPAREN, '`')
(OP, *)
(NUMBER,3)
```


## Scanner

- Lets write tokens for arithmetic expression:
(LPAREN, '(')
(NUMBER, 5)
(PLUS, +)
(NUMBER, 4)
(RPAREN, ' ')')
(TIMES, *)
(NUMBER, 3)
$(5+4) * 3$
(LPAREN, '(')
(FIVE, 5)
(PLUS, +)
(FOUR, 4)
(RPAREN, ' ')')
(TIMES, *)
(THREE, 3)


## Scanner

- Lets write tokens for arithmetic expression:
(LPAREN, '(')
(NUMBER, 5)
(PLUS, +)
(NUMBER, 4)
(RPAREN, ' ')')
(TIMES, *)
(NUMBER, 3)
(PAREN, '‘‘)
(NUMBER, 5)
(PLUS, +)
(NUMBER, 4)
(PAREN, ' ${ }^{\prime}$ ')
(TIMES, *)
(NUMBER, 3)

Defining tokens

## Defining tokens

- Literal - single character:
- PLUS = ' + ', TIMES = '*'


## Defining tokens

- Literal - single character:
- PLUS = '+', TIMES = ‘*'
- Keyword - single string:
- IF = "if", INT = "int"


## Defining tokens

- Literal - single character:
- PLUS = '+', TIMES = ‘*'
- Keyword - single string:
- IF = "if", INT = "int"
- Sets of words:
- NOUN = \{"Cat", "Dog", "Car"\}


## Defining tokens

- Literal - single character:
- PLUS = '+', TIMES = '*'
- Keyword - single string:
- IF = "if", INT = "int"
- Sets of words:
- NOUN = \{"Cat", "Dog", "Car"\}
- Numbers
- NUM = \{" 0 ", " 1 " ... $\}$


## Defining tokens

- Literal-single character:
-PLUS - ‘+', TIMES - ‘*'
- Keyword -single-string:
-IF = "if", INT = "int"
- Sets of words:
- NOUN - \{"Cat", "Dog", "Car" $\}$
- Numbers
- NUM - $\left\{{ }^{\prime \prime} 0 ", " 1 " \ldots\right\}$
- Regular expressions!


## Regular Expressions

- Lots of literature!
- Simplest grammar in the Chomsky language hierarchy
- abstract machine definition (finite automata)
- Many implementations (e.g. Python standard library)



## Regular Expressions

We will define RE's recursively:

The base case: a character literal

- The RE for a character ' $x$ ' is given by ' $x$ '. It matches only the character ' $x$ '

Examples: (demo)

## Regular Expressions

We will define RE's recursively:

Regular expressions are closed under concatenation:

- The concatenation of two REs $x$ and $y$ is given by $x y$ and matches the strings of RE $x$ concatenated with the strings of RE y

Examples (demo)

## Regular Expressions

We will define RE's recursively:

Regular expressions are closed under union:

- The union of two REs $x$ and $y$ is given by $x \mid y$ and matches the strings of RE $x$ or the strings of RE $y$

Examples (demo)

## Regular Expressions

We will define RE's recursively:

Regular expressions are closed under Kleene star:

- The Kleene star of an RE $x$ is given by $x^{*}$ and matches the strings of RE $x$ repeated 0 or more times

Examples (demo)

## Regular Expressions

- Use ()'s to force precedence!
- Without ()'s, what is the precedence of concatenation, union, and star?
-What are some experiments we can do?


## Regular Expressions

- Use ()'s to force precedence!
- Without ()'s, what is the precedence of concatenation, union, and star?
- star $>$ concatenation $>$ union


## Regular Expressions

Most RE implementations provide syntactic sugar:

- Ranges:
- [0-9]: any number between 0 and 9
- [a-z]: any lower case character
- [A-Z]: any upper case character
- Optional(?)
- Matches 0 or 1 instances:
- ab?c matches "abc" or "ac"
- can be implemented as: (abc | ac)


## Defining tokens using REs

- Literal - single character:
- PLUS = '+', TIMES = ‘*'
- Keyword - single string:
- IF = "if", INT = "int"
- Sets of words:
- NOUN = "(Cat)|(Dog)|(Car)"
- Numbers
- SINGLE_NUM $=[0-9]$

What about C-style IDs?

- NUM $=(-I \backslash+) ?[0-9]+(\backslash \cdot[0-9]+)$ ?


## Scanner Questions?

- A scanner splits a string into lexemes
- Tokens are defined using regular expressions
- Regular expressions are good for matching operators, parenthesis, variable names, numbers, key words etc.

Parser


## Parser

- Sentence:
- ARTICLE ADJECTIVE? NOUN VERB
-What about a mathematical sentence (expression)?


## Parser

- Sentence:
- ARTICLE ADJECTIVE? NOUN VERB
-What about a mathematical sentence (expression)?
- NUM


## Parser

- Sentence:
- ARTICLE ADJECTIVE? NOUN VERB
-What about a mathematical sentence (expression)?
- NUM
- NUM PLUS NUM


## Parser

- Sentence:
- ARTICLE ADJECTIVE? NOUN VERB
-What about a mathematical sentence (expression)?
- NUM
- NUM BIN_OP NUM


## Parser

- Sentence:
- ARTICLE ADJECTIVE? NOUN VERB
-What about a mathematical sentence (expression)?
- NUM
- NUM PLUS NUM
- NUM TIMES NUM
- NUM PLUS NUM TIMES NUM
- NUM PLUS NUM TIMES NUM
- NUM (BIN_OP NUM)*


## Context Free Grammars

- Backus-Naur form (BNF)

- A syntax for representing context free grammars
- Naturally create tree like structures
- More powerful than regular expressions


## Parser

- <production name> : <token>*
- Example: sentence: ARTICLE NOUN VERB
- <production name> : <token>* | <token>*
- Example:
sentence: ARTICLE ADJECTIVE NOUN VERB
| ARTICLE NOUN VERB


## Parser

- Production rules can reference other production rules
sentence: adjective_sentence
| non_adjective_sentence
adjective_sentence: ARTICLE ADJECTIVE NOUN VERB
non_adjective_sentence: ARTICLE NOUN VERB


## Parser

- Production rules can be recursive
- Imagine a list of adjectives:
"The small brown energetic dog barked"
sentence: ARTICLE adjective_list NOUN VERB


## Parser

- Production rules can be recursive
- Imagine a list of adjectives:
"The small brown energetic dog barked"
sentence: ARTICLE adjective_list NOUN VERB
adjective_list: ADJECTIVE adjective_list
| <empty>


## Next week

- Production rules for expressions
- parse trees
- associativity
- ambiguous grammars
- Homework is released next class:
- Have a look, but we will cover PLY and parsing with derivatives next Tuesday
- See you on Thursday!

