Name: $\qquad$ Login: $\qquad$ 1

## 1. Keys vs hash index: ( $\mathbf{1 0}$ points:)

Yes/No: In a hash table with chained elements for handling collisions, do elements in a chain have the same key?
No

## 2. Hashing: (40 points:)

Demonstrate what happens when we insert the keys $5,28,19,15,20,33,12,17$, 10, 11 into a hash table with collisions resolved by chaining. Let the table have 9 slots, and let the hash function be $h(k)=k \bmod 9$. Draw the contents of the hash table after all the keys have been inserted.

| k | $\mathrm{h}(\mathrm{k})$ |
| :---: | :---: |
| 5 | 5 |
| 28 | 1 |
| 19 | 1 |
| 15 | 6 |
| 20 | 2 |
| 33 | 6 |
| 12 | 3 |
| 17 | 8 |
| 10 | 1 |
| 11 | 2 |


| Hash table |  |
| :---: | :---: |
| key | value |
| 0 |  |
| 1 | $10,19,28$ |
| 2 | 11,20 |
| 3 | 12 |
| 4 |  |
| 5 | 5 |
| 6 | 33,15 |
| 7 |  |
| 8 | 17 |

$\qquad$

## 3. Open Addressing: (40 points:)

Consider inserting the keys $10,22,31,4,15,28,17,56,87,16$ into a hash table of length $\mathrm{m}=11$ using open addressing with the auxiliary hash function $h^{\prime}(\mathrm{k})=$ k. Show the contents of the hash table using quadratic probing with $\mathrm{c} 1=1$ and $\mathrm{c} 2=3$. Recall that $\mathrm{h}(\mathrm{k}, \mathrm{i})=\left(\mathrm{h}^{\prime}(\mathrm{k})+\mathrm{c} 1^{*} \mathrm{i}+\mathrm{c} 2 * \mathrm{i}^{2}\right) \bmod \mathrm{m}$.

| $\mathrm{h}(\mathrm{k}, \mathrm{i})$ | = | $\left(\mathrm{k}+\mathrm{i}+3 \mathrm{i}^{2}\right)$ |  | mod |
| :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{i}=0$ | $\mathrm{i}=1$ | $\mathrm{i}=2$ | $\mathrm{i}=3$ |
| k | h(k,i) |  |  |  |
| 10 | 10 |  |  |  |
| 22 | 0 |  |  |  |
| 31 | 9 |  |  |  |
| 4 | 4 |  |  |  |
| 15 | 4 | 8 |  |  |
| 28 | 6 |  |  |  |
| 17 | 6 | 10 | 9 | 3 |
| 56 | 1 |  |  |  |
| 87 | 10 | 3 | 2 |  |
| 16 | 5 |  |  |  |
|  |  |  |  |  |

11

| Hash table |  |
| :---: | :---: |
| key | value |
| 0 | 22 |
| 1 | 56 |
| 2 | 87 |
| 3 | 17 |
| 4 | 4 |
| 5 | 16 |
| 6 | 28 |
| 7 |  |
| 8 | 15 |
| 9 | 31 |
| 10 | 10 |

## 4. Open Ended Question: (10 points:)

The social security number in the US consists of 9 digits and are unique keys assigned to individuals, and are not re-used even when the person holding a number dies. Of the billion minus 1 possible keys, around 450 million have already been assigned. Associated with each social security number is information about us e.g. name, birthday, address, etc.
(a) (5 points:) Because social security numbers are (suppose to) uniquely identify an individual, they make excellent keys. If information about us are stored in a dictionary as key:value pairs in a hash table with chaining, how many slots should the hash table have if we are willing to search through the chain an average of 100 elements. That is a load factor of 100 .
$\mathrm{m}=450 \mathrm{M} / 100=4.5 \mathrm{M}$
(b) (5 points:) With recent data breaches and potential for identity thefts, many are requesting new social security numbers. That is, while a social security number uniquely identifies a person, a person can have more than one social security number. In a recent study conducted by ID Analytics, they found that around 20 million Americans have multiple social security numbers for various reasons - e.g. due to identity theft a person may request a new social security number. When issuing a new number to a person, the Social Security Administration will cross reference the two numbers to point to the same person.
Explain how you would handle such cases in hash tables with chaining?
add a link field to connect 2 nodes even if they're in different slots in hash table. downside is this will require space for another 450 M pointers even if just a few are needed. do accept other solutions if they make sense even if not efficient.

