

Extending the TPTP Language to Higher-Order Logic with Automated Parser Generation

System Description

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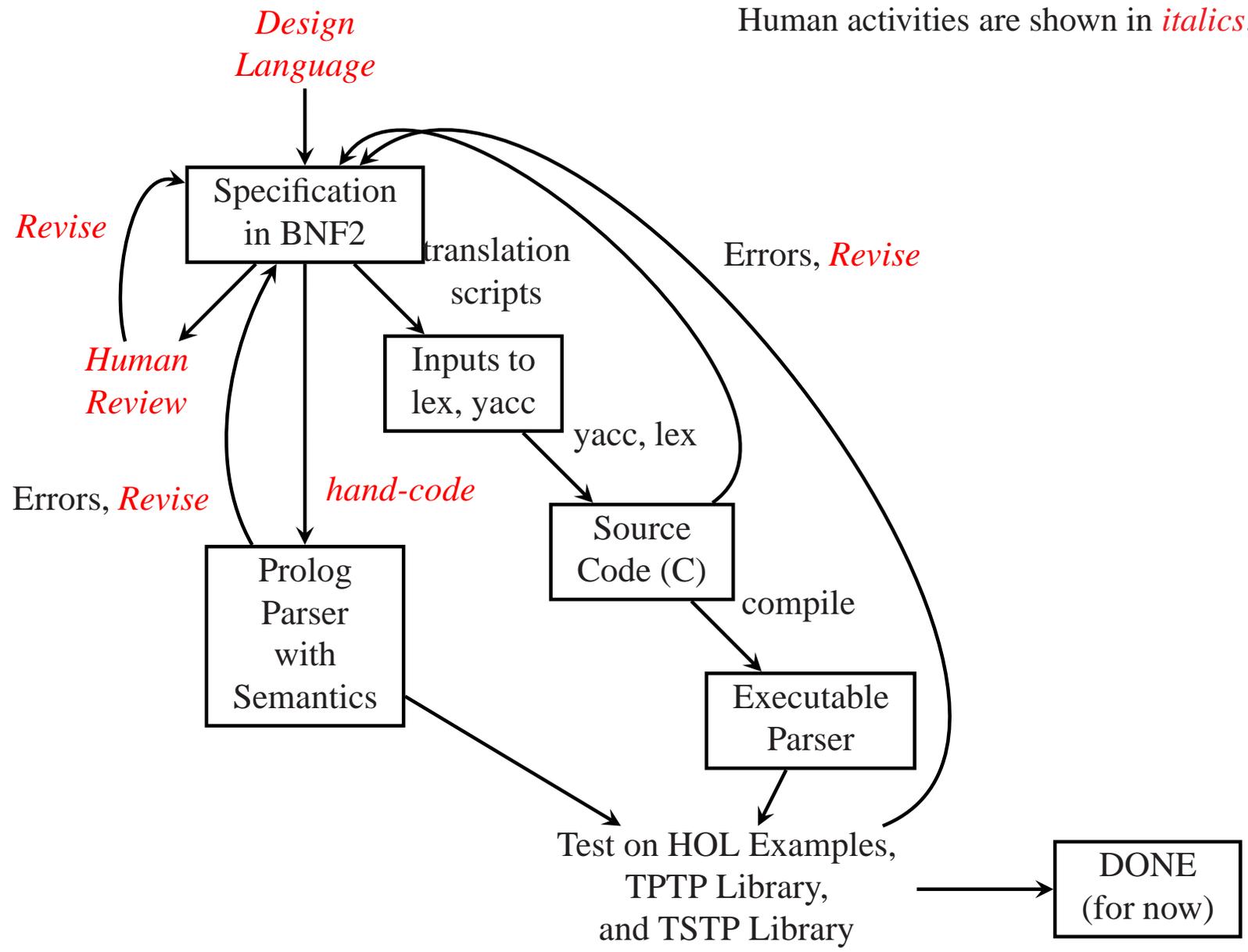
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`http://www.cse.ucsc.edu/~avg/TPTPparser/`

`http://www.cs.miami.edu/~tptp/`

Overview of the System to Evaluate a Proposed HOTPTP Language

Human activities are shown in *italics*.



BNF2 – A Two-Level Specification Language

Traditional BNF (Backus-Naur Form) has one syntax for all grammar symbols.

Modern practice uses two levels of symbols

- *tokens* are defined by regular expressions, and are treated as terminal symbols in the context-free part of the grammar.
- *grammar symbols* are the nonterminal symbols of the context-free grammar.

<i>Meta-Symbol</i>	<i>Rule Type</i>	<i>Examples</i> (some are simplified from the TPTP language)	
<code>::=</code>	Grammar	<code><TPTP input></code>	<code>::= <annotated formula> <comment></code> <code><nonassoc op></code> <code>::= <=> => <= <~></code> <code><formula role></code> <code>::= <lower word></code>
<code>::=</code>	Semantic	<code><formula role></code>	<code>::= axiom conjecture lemma theorem negated_conjecture</code>
<code>::-</code>	Token	<code><lower word></code>	<code>::- <lower><alphanum>*</code>
<code>::=</code>	Macro	<code><lower></code> <code><alphanum></code>	<code>::: [a-z]</code> <code>::: [A-Za-z0-9_]</code>

Proposed HOTPTP Syntax, Extending TPTP to Higher-Order Logic

New operators continue TPTP tradition that operators are logical symbols:

- Lambda abstraction: λ (synonym `lambda`)
- Function application: `@`
- “Maps to” for types: \rightarrow (synonym $>$)
- “Defined as”: `:=` (Hudak [1989] uses \equiv which looks reflexive.)

Colon operator (`:`) has several new meanings.

```
hof(1, definition,  
    set_union := lambda [A: $type]: lambda [D: ((A-> $o)-> $o), X: A]:  
                ? [S: (A-> $o)]: ( (D @ S) & (S @ X) ) ).
```

The lambda expression shown is:

$$\lambda A:\tau. \lambda D:((A \rightarrow o) \rightarrow o). \lambda X:A. \exists S:(A \rightarrow o). ((D @ S) \wedge (S @ X)).$$

More details and examples at <http://www.cse.ucsc.edu/~avg/TPTPparser/>.

To be discussed at TPTP Tea Party.